

IN THE CLAIMS:

1. (Currently Amended) A method for operating a data storage system, comprising:

- creating a writable virtual disk (vdisk) at a selected time, the writable vdisk referencing changes in data stored in the data storage system after the writable vdisk was created;
- maintaining a backing store, the backing store referencing data stored in the data storage system which has not been changed since the writable vdisk was created;
- loading blocks of the writable vdisk into a memory, the loaded blocks including a writable vdisk indirect block having a plurality of fields, each field storing a valid pointer to a data block or an invalid pointer representing ~~a particular hole~~ one of a plurality of holes, where each hole instructs the data storage system to examine a corresponding virtual block number pointer in the backing store;
- loading blocks of the backing store into the memory, the loaded blocks including a backing store indirect block having a plurality of fields, each backing store indirect block field corresponding to a field of the writable vdisk indirect block, one or more backing store indirect block fields having a pointer to a data block;
- searching each field of the writable vdisk indirect block for a hole; and
- filling each hole in the writeable vdisk by replacing each invalid pointer field ~~having a hole in the writeable vdisk indirect block~~ with a new pointer to the data block referenced by the corresponding backing store indirect block field to update the writable vdisk to reference both the data which is unchanged since the writable vdisk was created and the data which has been changed since the writable vdisk was created.

2. (Previously Presented) The method of claim 1, further comprising:

- dirtying the data block pointed to by the backing store indirect block to enable write allocation of the dirty data block without altering a data content of the data block.

3. (Previously Presented) The method of claim 1, further comprising:

2 choosing a new pointer for a newly allocated data block containing an unaltered
3 data content;
4 setting bits in block allocation structures for the newly allocated data block; and
5 placing the new pointer to the newly allocated data block into the field of the
6 writable vdisk indirect block to replace the hole.

1 4. (Previously Presented) The method of claim 3 further comprising:

2 freeing the dirty data block; and
3 writing the newly allocated data block to disk.

1 5. (Previously Presented) The method of claim 4 further comprising:

2 releasing an association of the writable vdisk to the backing store to thereby
3 separate the writable vdisk data blocks from the backing store data blocks.

1 6. (Original) The method of claim 1 wherein the pointers contained in the writable vdisk
2 indirect block fields and the backing store indirect block fields comprise logical volume
3 block numbers (VBNs).

1 7. (Original) The method of claim 1 wherein the invalid pointers contained in the
2 writable vdisk indirect block fields comprise a zero logical volume block number (VBN).

1 8. (Original) The method of claim 1 wherein the plurality of fields in the writable vdisk
2 indirect block are a writable vdisk level 1 buffer and the plurality of fields in the backing
3 store indirect block are a backing store level 1 buffer.

1 9. (Currently Amended) An apparatus for operating a computer database, comprising:
2 a writable virtual disk (vdisk) created at a selected time, the writable vdisk
3 referencing changes in data stored in a data storage system after the writable vdisk was
4 created;

a backing store, the backing store referencing data stored in the data storage system which has not been changed since the writable vdisk was created;

a backdoor message handler adapted to load blocks of the writable vdisk and backing store from disk into a memory of the storage system;

a writable vdisk indirect block in the memory having a plurality of fields, each field storing a valid pointer to a data block or an invalid pointer representing a particular ~~hole~~ of a plurality of holes, where each hole instructs the data storage system to examine a corresponding virtual block number pointer in the backing store;

a backing store indirect block in the memory having a plurality of fields, each backing store indirect block field corresponding to a field of the writable vdisk indirect block, each backing store indirect block field having a pointer to a data block;

a special loading function for searching each field of the writable vdisk indirect block for one or more fields representing a hole; and

a write allocator for filling each hole in the writeable vdisk by replacing each invalid pointer field representing a hole in the writable vdisk indirect block with a new pointer to the data referenced by the corresponding backing store indirect block field to update the writable vdisk to reference both the data which is unchanged since the writable vdisk was created and the data which has been changed since the writable vdisk was created.

10. (Previously Presented) The apparatus of claim 9 wherein the write allocator further comprises:

a new pointer for a newly allocated data block containing an unaltered data content, set bits in block allocation structures for the newly allocated data block, and place the new pointer to the newly allocated data block into the field of the writable vdisk indirect block to replace the hole.

11. (Original) The apparatus of claim 10 wherein the write allocator is further adapted to:

free the dirty data block and write the newly allocated data block to disk.

1 12. (Original) The apparatus of claim 9 wherein the backdoor message handler loads the
2 blocks of the writable vdisk and the blocks of the backing store during periods of reduced
3 processing activity.

1 13. (Original) The apparatus of claim 9 wherein the pointers contained in the writable
2 vdisk indirect block fields and the backing store indirect block fields comprise logical
3 volume block numbers (VBNs).

1 14. (Original) The apparatus of claim 9 wherein the invalid pointers contained in the
2 writable vdisk indirect block fields comprise a zero logical volume block number (VBN).

1 15. (Original) The apparatus of claim 9 wherein the plurality of fields in the writable
2 vdisk indirect block comprises a writable vdisk level 1 buffer and the plurality of fields in
3 the backing store indirect block comprises a backing store level 1 buffer.

1 16.-18. (Cancelled).

1 19. (Currently Amended) A data storage system apparatus, comprising:
2 means for creating a writable virtual disk (vdisk) at a selected time, the writable
3 vdisk referencing changes in data stored in the data storage system after the writable
4 vdisk was created;
5 means for maintaining a backing store, the backing store referencing data stored
6 in the data storage system which has not been changed since the writable vdisk was
7 created;
8 means for loading blocks of the writable vdisk from a disk into a memory, the
9 loaded blocks including a writable vdisk indirect block having a plurality of fields, each
10 field storing a valid pointer to a data block or an invalid pointer representing a particular
11 ~~hole~~ hole of a plurality of holes, where each hole instructs the data storage system to
12 examine a corresponding virtual block number pointer in the backing store;

means for loading blocks of the backing store from a disk into the memory, the loaded blocks including a backing store indirect block having a plurality of fields, each backing store indirect block field corresponding to a field of the writable vdisk indirect block, one or more backing store indirect block fields having a pointer to a data block;

means for searching each field of the writable vdisk indirect block for a hole; and

means for filling each hole in the writable vdisk by replacing each invalid pointer field having a hole in the writable vdisk indirect block with a new pointer to the data block referenced by the corresponding backing store indirect block field to update the writable vdisk to reference both the data which is unchanged since the writable vdisk was created and the data which has been changed since the writable vdisk was created .

20. (Currently Amended) A computer readable medium, ~~including program instructions executing on a computer, executable~~ the program instructions executed by a processor, comprising ~~including instructions for performing the steps of:~~

creating program instructions that create a writable virtual disk (vdisk) at a selected time, the writable vdisk referencing changes in data stored in a data storage system after the writable vdisk was created;

maintaining program instructions that maintain a backing store, the backing store referencing data stored in the data storage system which has not been changed since the writable vdisk was created;

loading program instructions that load blocks of the writable vdisk from a disk into a memory, the loaded blocks including a writable vdisk indirect block having a plurality of fields, each field storing a valid pointer to a data block or an invalid pointer representing ~~a particular hole~~ one of a plurality of holes, where each hole instructs the data storage system to examine a corresponding virtual block number pointer in the backing store;

loading program instructions that load blocks of the backing store from a disk into the memory, the loaded blocks including a backing store indirect block having a plurality of fields, each backing store indirect block field corresponding to a field of the writable

vdisk indirect block, one or more backing store indirect block fields having a pointer to a data block;

~~searching program instructions that search~~ each field of the writable vdisk indirect block for a hole; and

~~program instructions that fill each hole in the writeable vdisk by replacing each invalid pointer field having a hole in the writable vdisk indirect block with a new pointer~~ to the data block referenced by the corresponding backing store indirect block field to update the writable vdisk to reference both the data which is unchanged since the writable vdisk was created and the data which has been changed since the writable vdisk was created.

21-22. (Cancelled).

23. (Currently Amended) A method for operating a data storage system, comprising:

creating a writable virtual disk (vdisk) at a selected time, the writable vdisk referencing changes in data stored in the data storage system after the writable vdisk was created, the writable vdisk having a plurality of holes where each hole instructs the storage system to examine a corresponding virtual block number pointer in a backing store;

maintaining the backing store, the backing store referencing the data stored in the data storage system which has not been changed since the writable vdisk was created;

searching each field of the writable vdisk for a hole; and

~~filling each hole in the writable vdisk by replacing~~ filling each hole in the writeable vdisk by replacing referencing each hole in the writable vdisk to point to the data block referenced by the corresponding backing store indirect block to effectively fill each hole of the writeable vdisk with the data block and thus update the writable vdisk to reference both the data which is unchanged since the writable vdisk was created and the data which has been changed since the writable vdisk was created.

24. (Previously Presented) The method of claim 23, further comprising:

dirtying the data block pointed to by the backing store indirect block to enable
write allocation of the dirty data block without altering a data content of the data block.

25. (Previously Presented) The method of claim 23 further comprising:
choosing a new pointer for a newly allocated data block containing an unaltered
data content;
setting bits in block allocation structures for the newly allocated data block; and
placing the new pointer to the newly allocated data block into the field of the
writable vdisk indirect block to replace the hole.

26. (Previously Presented) The method of claim 25, further comprising:
freeing the dirty data block; and
writing the newly allocated data block to disk.

27. (Previously Presented) The method of claim 26 further comprising:
releasing an association of the writable vdisk to the backing store to thereby
separate the writable vdisk data blocks from the backing store data blocks.

28. (Previously Presented) The method of claim 23, further comprising:
including logical volume block numbers (VBNs) in the pointers contained in the
writable vdisk indirect block fields and the backing store indirect block fields.

29. (Previously Presented) The method of claim 23, further comprising:
using a zero logical volume block number (VBN) as the invalid pointers
contained in the writable vdisk indirect block fields.

30. (Previously Presented) The method of claim 23, further comprising:
using a writable vdisk level 1 buffer for the plurality of fields in the writable vdisk
indirect block and using a backing store level 1 buffer for the plurality of fields in the
backing store indirect block.

1 31. (Currently Amended) A data storage system, comprising:

2 a writable virtual disk (vdisk) created at a selected time, the writable vdisk
3 referencing changes in data stored in the data storage system after the writable vdisk was
4 created, the writable vdisk having a plurality of holes, each hole instructing the storage
5 system to examine a corresponding virtual block number pointer in a backing store;

6 the backing store referencing data stored in the data storage system which has not
7 been changed since the writable vdisk was created;

8 a processor to search each field of the writable vdisk for a hole; and

9 the processor to fill each hole in the writable vdisk so reference that each hole in
10 the writable vdisk to point points to the data block referenced by the corresponding
11 backing store indirect block to effectively fill each hole of the writeable vdisk with the
12 data block and thus update the writable vdisk to reference both the data which is
13 unchanged since the writable vdisk was created and the data which has been changed
14 since the writable vdisk was created.

1 32. (Previously Presented) The system of claim 31, further comprising:

2 the data block pointed to by the backing store indirect block are dirtied to enable
3 write allocation of the dirty data block without altering a data content of the data block.

1 33. (Previously Presented) The system of claim 31 further comprising:

2 a new pointer chosen for a newly allocated data block containing an unaltered
3 data content;

4 bits are set in a block allocation structures for the newly allocated data block; and

5 a new pointer to the newly allocated data block placed into a field of the writable
6 vdisk indirect block to replace the hole.

1 34. (Previously Presented) The system of claim 33, further comprising:

2 the dirty data block is freed; and

3 the newly allocated data block is written to disk.

1 35. (Previously Presented) The system of claim 34 further comprising:
2 an association of the writable vdisk to the backing store is released to thereby
3 separate the writable vdisk data blocks from the backing store data blocks.

1 36. (Previously Presented) The system of claim 31, further comprising:
2 logical volume block numbers (VBNs) included in the pointers contained in the
3 writable vdisk indirect block fields and the backing store indirect block fields.

1 37. (Previously Presented) The system of claim 31, further comprising:
2 a zero logical volume block number (VBN) used as the invalid pointers contained
3 in the writable vdisk indirect block fields.

1 38. (Previously Presented) The system of claim 31, further comprising:
2 a writable vdisk level 1 buffer used for the plurality of fields in the writable vdisk
3 indirect block and a backing store level 1 buffer used for the plurality of fields in the
4 backing store indirect block.

1 39. (Currently Amended) A computer readable ~~medium~~ comprising:
2 ~~said computer readable media~~ containing executable program instructions
3 executed for execution only a processor ~~for a method of method for operating a data~~
4 ~~storage system, the method having comprising:~~
5 creating program instructions that create a writable virtual disk (vdisk) at a
6 selected time, the writable vdisk referencing changes in data stored in the data storage
7 system after the writable vdisk was created, the writable vdisk having a plurality of holes
8 where each hole instructs the storage system to examine a corresponding virtual block
9 number pointer in a backing store;
10 maintaining program instructions that maintain the backing store, the backing
11 store referencing data stored in the data storage system which has not been changed since
12 the writable vdisk was created;

13 | ~~searching-program instructions that search~~ each field of the writable vdisk for a
14 | hole; and
15 | ~~referencing-program instructions that fill~~ each hole in the writable vdisk to point
16 | to the data block referenced by the corresponding backing store indirect block to
17 | ~~effectively fill each hole of the writeable vdisk with the data block and thus~~ update the
18 | writable vdisk to reference both the data which is unchanged since the writable vdisk was
19 | created and the data which has been changed since the writable vdisk was created.

1 | 40. (Currently Amended) A method for operating a data storage system, comprising:
2 | creating a writable virtual disk (vdisk) at a selected time, the writable vdisk
3 | referencing changes in data stored in the data storage system after the writable vdisk was
4 | created, the writable vdisk having a plurality of holes where each hole instructs the data
5 | storage system to examine a corresponding virtual block number pointer in a backing
6 | store;
7 | maintaining the backing store, the backing store referencing the data stored in the
8 | data storage system which has not been changed since the writable vdisk was created;
9 | searching, by a background task process, each field of the writable vdisk for a
10 | hole; ~~and~~
11 | for each hole in the writeable vdisk, marking as dirty the corresponding data block
12 | pointed to by the backing store indirect block without modifying the corresponding data
13 | block; and
14 | performing a write allocation to replace ~~referencing~~ each hole in the writable
15 | vdisk to point to the data block marked as dirty and referenced by the corresponding
16 | backing store indirect block to update the writable vdisk to reference both the data which
17 | is unchanged since the writable vdisk was created and the data which has been changed
18 | since the writable vdisk was created.

1 | 41. (Currently Amended) A data storage system, comprising:
2 | a writable virtual disk (vdisk) created at a selected time, the writable vdisk
3 | referencing changes in data stored in the data storage system after the writable vdisk was

created, the writable vdisk having a plurality of holes where each hole instructs the data storage system to examine a corresponding virtual block number pointer in the backing store;

the backing store referencing the data stored in the data storage system which has not been changed since the writable vdisk was created;

a background task processor to search each field of the writable vdisk for a hole; and

the background task processor to mark as dirty, for each hole in the writeable vdisk, the corresponding data block pointed to by the backing store indirect block without modifying the corresponding data block, and perform a write allocation to ~~replacereference~~ each hole in the writable vdisk to point to the data block marked as dirty ~~and~~ referenced by the corresponding backing store indirect block to update the writable vdisk to reference both the data which is unchanged since the writable vdisk was created and the data which has been changed since the writable vdisk was created.

42. (Currently Amended) A computer readable media, ~~comprising:~~

~~—said computer readable media containing executable program instructions executed by for execution on a processor, comprising: for a method of method for operating a data storage system, the method having,~~

creating program instructions that create a writable virtual disk (vdisk) at a selected time, the writable vdisk referencing changes in data stored in the data storage system after the writable vdisk was created, the writable vdisk having a plurality of holes where each hole instructs the data storage system to examine a corresponding virtual block number pointer in a backing store;

maintaining program instructions that maintain the backing store, the backing store referencing the data stored in the data storage system which has not been changed since the writable vdisk was created;

searching program instructions that search, by a background task process, each field of the writable vdisk for a hole; ~~and~~

15 program instructions that mark as dirty, for each hole in the writeable vdisk, the
16 corresponding data block pointed to by the backing store indirect block without
17 modifying the corresponding data block; and

18 program instructions that perform a write allocation to replace ~~referencing~~ each
19 hole in the writable vdisk to point to the data block marked as dirty and referenced by the
20 corresponding backing store indirect block to update the writable vdisk to reference both
21 the data which is unchanged since the writable vdisk was created and the data which has
22 been changed since the writable vdisk was created.